

# Everything is Vecchia: Unifying low-rank and sparse inverse Cholesky approximations

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# Framing the project

**Past work:** Randomized numerical linear algebra (RNLA) has generated low-rank approximations with strong theoretical guarantees:

- Randomized SVD
- Generalized Nystrom
- Randomized block Krylov.

**Problem:** What if the matrix is not approximately low-rank? How do we solve linear algebra problems then?

**Goal:** We need a flexible class of matrix approximations that includes low-rank approximation and admits strong guarantees.

**Partial progress:** The Vecchia approximation is a solid option. It includes low-rank approximations and other *mysterious* structures. We should understand it better.

**Outlook:** The Vecchia approximation will not solve positive-semidefinite approximation, but it will take us closer.

[Joint work with Eagan Kaminetz, arXiv paper coming out next week.]

# Cholesky decomposition

Let us focus on positive-semidefinite matrices  $\mathbf{A} \in \mathbb{C}^{n \times n}$ .

They can be represented by a Cholesky or inverse Cholesky decomposition

$$\mathbf{A} = \mathbf{P}\mathbf{L}\mathbf{D}\mathbf{L}^*\mathbf{P}^* \quad \text{or} \quad \mathbf{A} = \mathbf{P}\mathbf{C}^{-1}\mathbf{D}\mathbf{C}^{-*}\mathbf{P}^*.$$

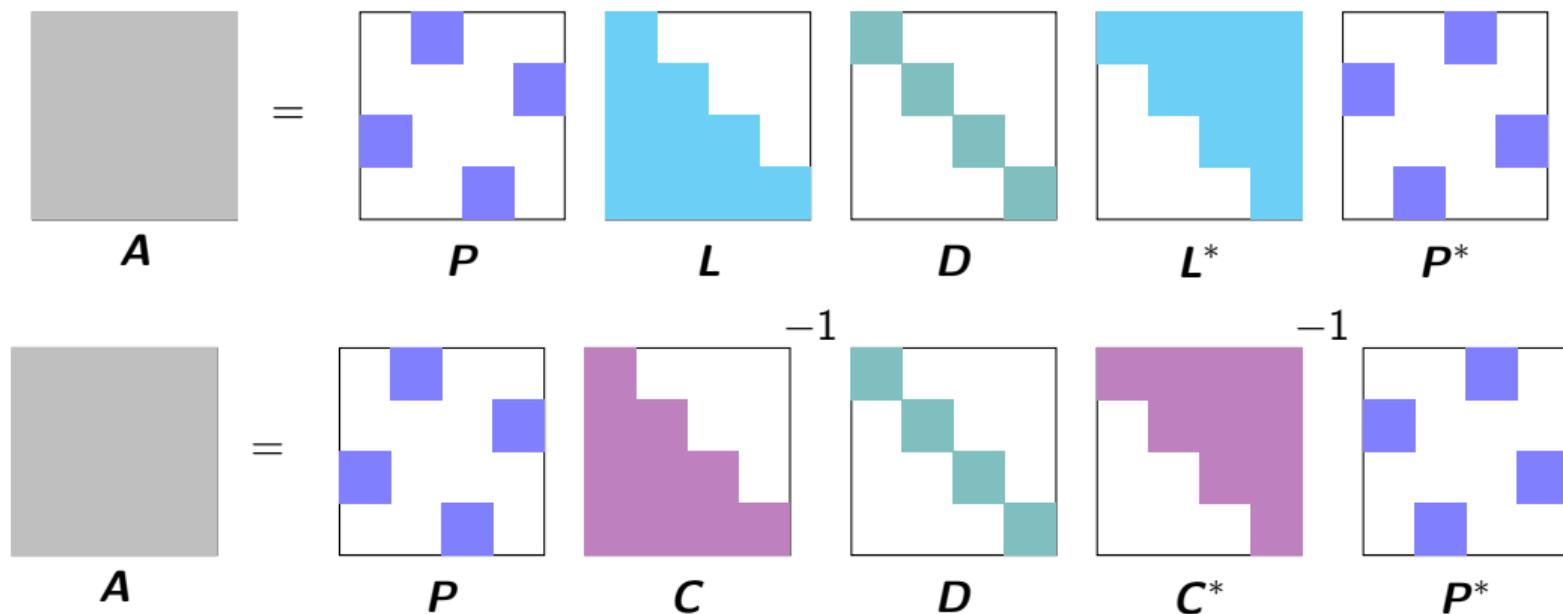
Here,  $\mathbf{P} \in \{0, 1\}^{n \times n}$  is a permutation matrix,  $\mathbf{L}$  is lower triangular with ones on the diagonal, and  $\mathbf{D}$  is diagonal with nonnegative entries.

These decompositions exist for every positive-semidefinite matrix  $\mathbf{A}$  and permutation matrix  $\mathbf{P}$ , they are equivalent using  $\mathbf{L} = \mathbf{C}^{-1}$ .

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Given such a factorization, we can efficiently solve linear systems, compute determinants, etc. in  $\mathcal{O}(n^2)$  or  $\mathcal{O}(n)$  time.

# Cholesky decomposition



**Figure:** Cholesky and inverse Cholesky decompositions of a dense matrix  $A$ . Factors  $P, L, D$  or  $P, C, D$  are stored, and inverses  $C^{-1}, C^{*-1}$  are accessed implicitly. Filled boxes show entries that are allowed to be nonzero.

# Sparse Cholesky approximation

Recent work uses *sparse* Cholesky or inverse Cholesky approximations

$$\hat{\mathbf{A}} = \mathbf{P}\hat{\mathbf{L}}\hat{\mathbf{D}}\hat{\mathbf{L}}^* \mathbf{P}^* \quad \text{or} \quad \hat{\mathbf{A}} = \mathbf{P}\hat{\mathbf{C}}^{-1}\hat{\mathbf{D}}\hat{\mathbf{C}}^{-*} \mathbf{P}^*.$$

Here,  $\hat{\mathbf{L}}$  or  $\hat{\mathbf{C}}$  is very sparse (but  $\hat{\mathbf{L}}^{-1}$  or  $\hat{\mathbf{C}}^{-1}$  may not be)

- “randomized Cholesky” (Kyng, Sachdeva),
  - “sparse Cholesky” (Schäfer, Katzfuss, Owhadi),
  - “randomly pivoted Cholesky” (Chen, Epperly, Tropp, Webber)
- 
- The sparsity pattern  $S_i \subseteq \{1, \dots, i-1\}$  lists the nonzero off-diagonal entries in row  $i$  of  $\hat{\mathbf{L}}$  or  $\hat{\mathbf{C}}$ .
  - If  $|S_i| \leq s$  for  $i = 1, \dots, n$ , we can solve linear systems in  $\mathcal{O}(ns)$  operations.
  - Often we can generate  $\hat{\mathbf{A}}$  in  $\mathcal{O}(s^2n)$  or  $\mathcal{O}(s^3n)$  operations, cheaper than the cost of looking at each entry once.

# Sparse Cholesky approximation

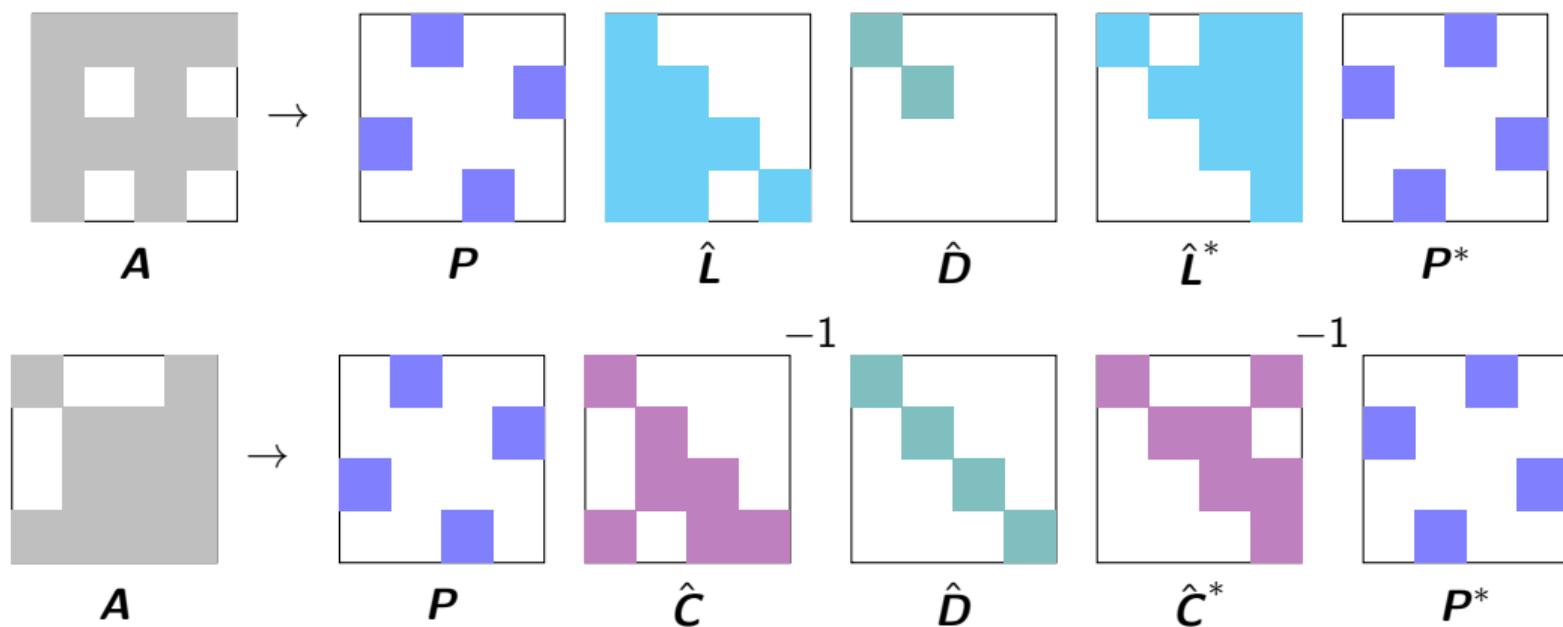
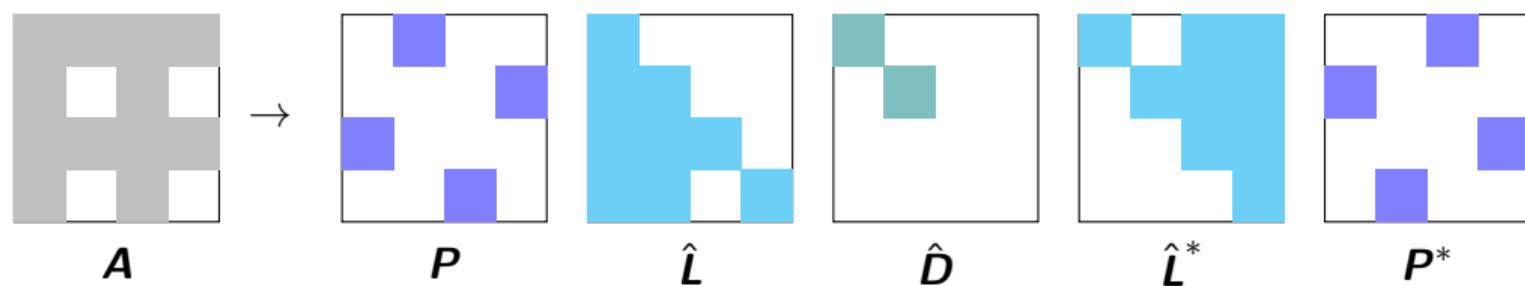


Figure: Sparse Cholesky and inverse Cholesky approximations of a dense matrix  $\mathbf{A}$ . Filled boxes show entries that are allowed to be nonzero.

# Partial pivoted Cholesky

*Partial pivoted Cholesky* generates a rank- $r$  sparse Cholesky approximation that matches user-selected columns, indexed by  $u_1, \dots, u_r$ .

Forming this approximation requires  $\mathcal{O}(rn)$  entry look-ups and  $\mathcal{O}(r^2n)$  extra processing.

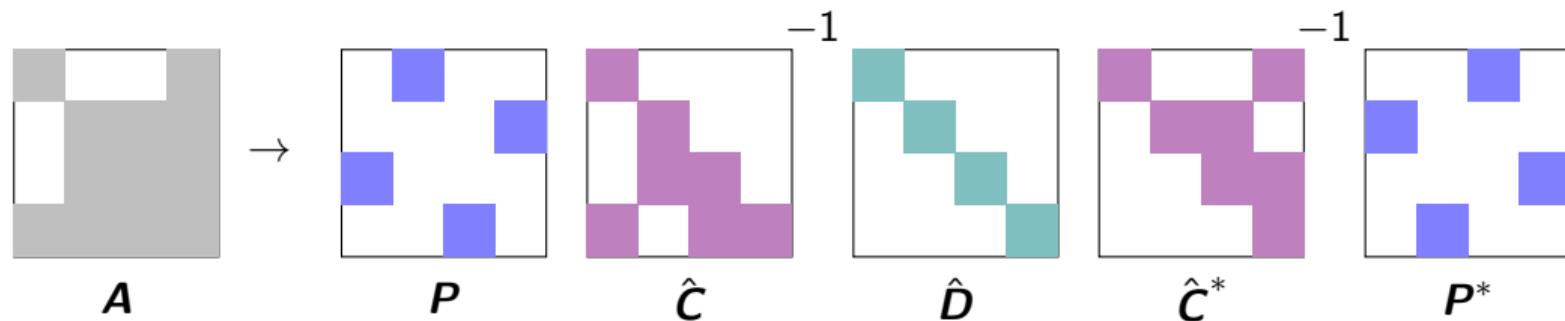


**Figure:** Partial pivoted Cholesky accesses gray-colored entries of  $\mathbf{A}$ . The rank is  $r = 2$ . Columns  $u_1 = 3$  and  $u_2 = 1$  are perfectly replicated.

# Vecchia

*Vecchia* approximation (named after Aldo Vecchia) generates a sparse inverse Cholesky approximation with any user-selected permutation  $\mathbf{P}$  and sparsity pattern  $\{S_i\}_{i=1}^r$ .

Traditionally requires  $\mathcal{O}(s^2n)$  entry look-ups and  $\mathcal{O}(s^3n)$  extra processing, where  $s$  is an upper bound on the cardinality,  $|S_i| \leq s$ .



**Figure:** Vecchia approximation accesses gray entries of  $\mathbf{A}$ . Pivots are  $u_1 = 3$ ,  $u_2 = 1$ ,  $u_3 = 4$ ,  $u_4 = 2$ . Sparsity pattern is  $S_2 = \emptyset$ ,  $S_3 = \{2\}$ ,  $S_4 = \{1, 3\}$ .

# Formula for the Vecchia approximation

Vecchia approximation uses  $\hat{\mathbf{C}}(i, S_i)$  and  $\hat{\mathbf{D}}(i, i)$  that solve

$$\begin{bmatrix} \hat{\mathbf{C}}(i, S_i) & 1 \end{bmatrix} \begin{bmatrix} \mathbf{A}(S_i, S_i) & \mathbf{A}(S_i, i) \\ \mathbf{A}(i, S_i) & \mathbf{A}(i, i) \end{bmatrix} = \begin{bmatrix} \mathbf{0} & \hat{\mathbf{D}}(i, i) \end{bmatrix}, \quad \text{for } i = 1, \dots, n.$$

# Partial Cholesky + Vecchia

Traditionally, partial Cholesky and Vecchia are used for different tasks:

- Partial Cholesky exposes **low-rank structure**.
- Vecchia exposes **sparse structure** in the inverse Cholesky factor.

Edmond Chow, Yuanzhe Xi, and colleagues (2024 SISC, 2025 SIMAX) combined the approximations, by forming a partial Cholesky approximation *first* and forming a residual Vecchia approximation *second*.

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The Chow & Xi papers were empirical, with promising numerical experiments. I wanted to find out more.

# Partial Cholesky + Vecchia

## Theorem (Partial Cholesky + Vecchia = Vecchia)

Given a target positive-semidefinite matrix  $\mathbf{A} \in \mathbb{C}^{n \times n}$ , consider the following two-part approximation.

1. Generate a partial Cholesky approximation of  $\mathbf{A}$  with permutation  $\mathbf{P}$  and approximation rank  $r$ . Call it  $\hat{\mathbf{A}}_{\text{part}}$ .
2. Generate a Vecchia approximation of the residual  $\mathbf{R} = \mathbf{A} - \hat{\mathbf{A}}_{\text{part}}$  with permutation  $\mathbf{P}$  and sparsity pattern  $\{Q_i\}_{i=1}^n$ . Call it  $\hat{\mathbf{A}}_{\text{res}}$ .

Then  $\hat{\mathbf{A}}_{\text{part}} + \hat{\mathbf{A}}_{\text{res}}$  can be rewritten as a Vecchia approximation of  $\mathbf{A}$  with permutation  $\mathbf{P}$  and an augmented sparsity pattern  $S_i = (\{1, \dots, r\} \cup Q_i) \cap \{1, \dots, i-1\}$ .

## Proof.

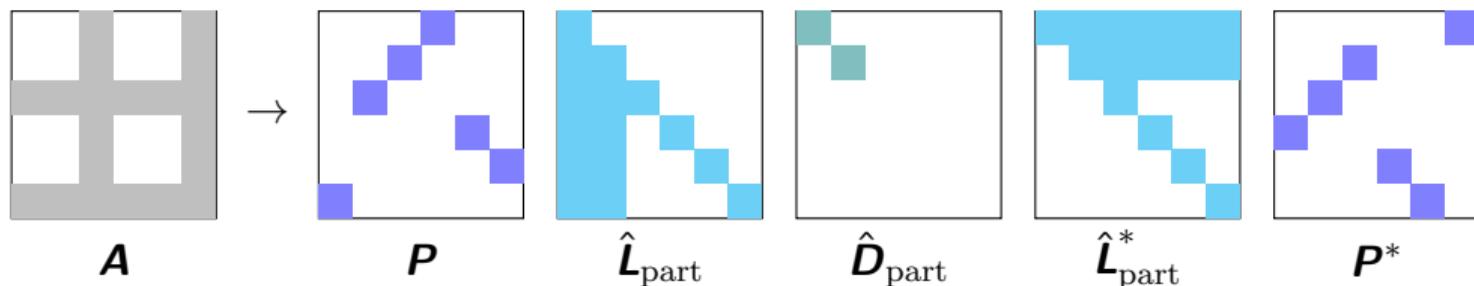
Proof in pictures (next three slides).



# Partial Cholesky + Vecchia

Partial Cholesky accesses  $\mathbf{A}$  and generates

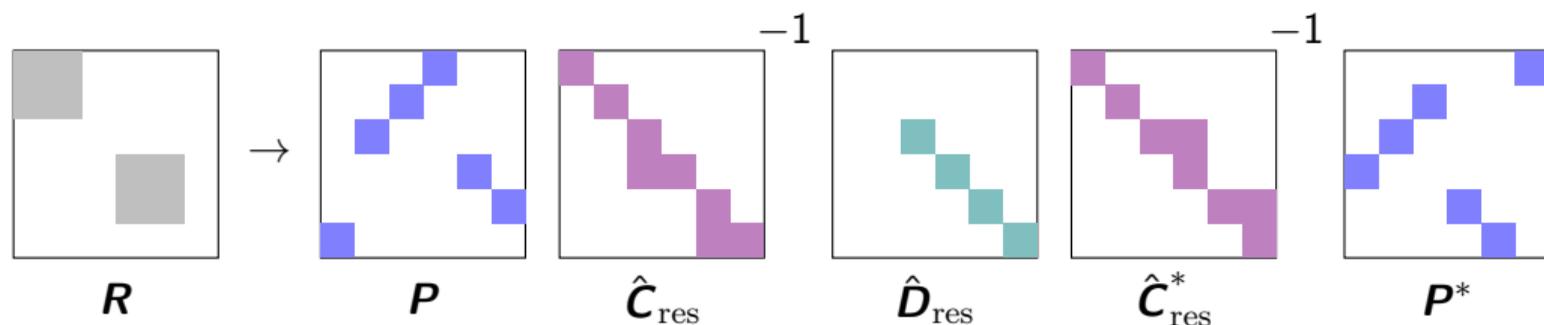
$$\hat{\mathbf{A}}_{\text{part}} = \mathbf{P} \hat{\mathbf{L}}_{\text{part}} \hat{\mathbf{D}}_{\text{part}} \hat{\mathbf{L}}_{\text{part}}^* \mathbf{P}^*.$$



# Partial Cholesky + Vecchia

Vecchia uses  $R = A - \hat{A}_{\text{part}}$  to generate

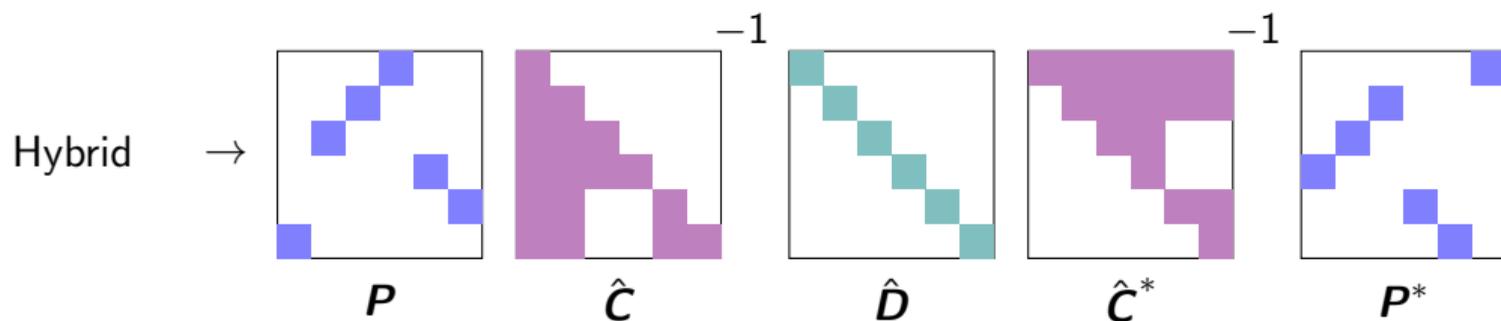
$$\hat{A}_{\text{res}} = P \hat{C}_{\text{res}}^{-1} \hat{D}_{\text{res}} \hat{C}_{\text{res}}^{-*} P^*.$$



# Partial Cholesky + Vecchia

Partial Cholesky + Vecchia can be rewritten

$$\begin{aligned}\hat{\mathbf{A}} &= \hat{\mathbf{A}}_{\text{part}} + \hat{\mathbf{A}}_{\text{res}} \\ &= \mathbf{P}\hat{\mathbf{C}}^{-1}\hat{\mathbf{D}}\hat{\mathbf{C}}^{-*}\mathbf{P}^*.\end{aligned}$$



This is a sparse inverse Cholesky decomposition with the right permutation and sparsity pattern. The rest of the proof just checks the Vecchia linear systems. ■

## Why is this important?

Partial Cholesky + Vecchia is often **computationally efficient**.

- + It reduces Vecchia's cost from  $\mathcal{O}(s^2n)$  entry look-ups and  $\mathcal{O}(s^3n)$  operations to  $\mathcal{O}(sn)$  look-ups and  $\mathcal{O}(s^2n)$  operations for a special sparsity pattern with  $|S_i| \leq s$ .

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Partial Cholesky + Vecchia is often **accurate**.

- + In the simplest case, we apply partial Cholesky to  $\mathbf{A}$  and then apply Vecchia with the minimal sparsity pattern  $S_i = \emptyset$ . This leads to partial Cholesky + diagonal

$$\hat{\mathbf{A}} = \hat{\mathbf{A}}_{\text{part}} + \text{diag}(\mathbf{A} - \hat{\mathbf{A}}_{\text{part}}).$$

- + Alternatively, we can apply partial Cholesky to  $\mathbf{A}$  and then apply Vecchia with a small, carefully chosen sparsity pattern.

# Experimental setup

We downloaded 27 machine learning data sets with 4–784 predictors and 1 quantitative response (OpenML, LibSVM, ...). We subsampled  $n = 20,000$  data points.

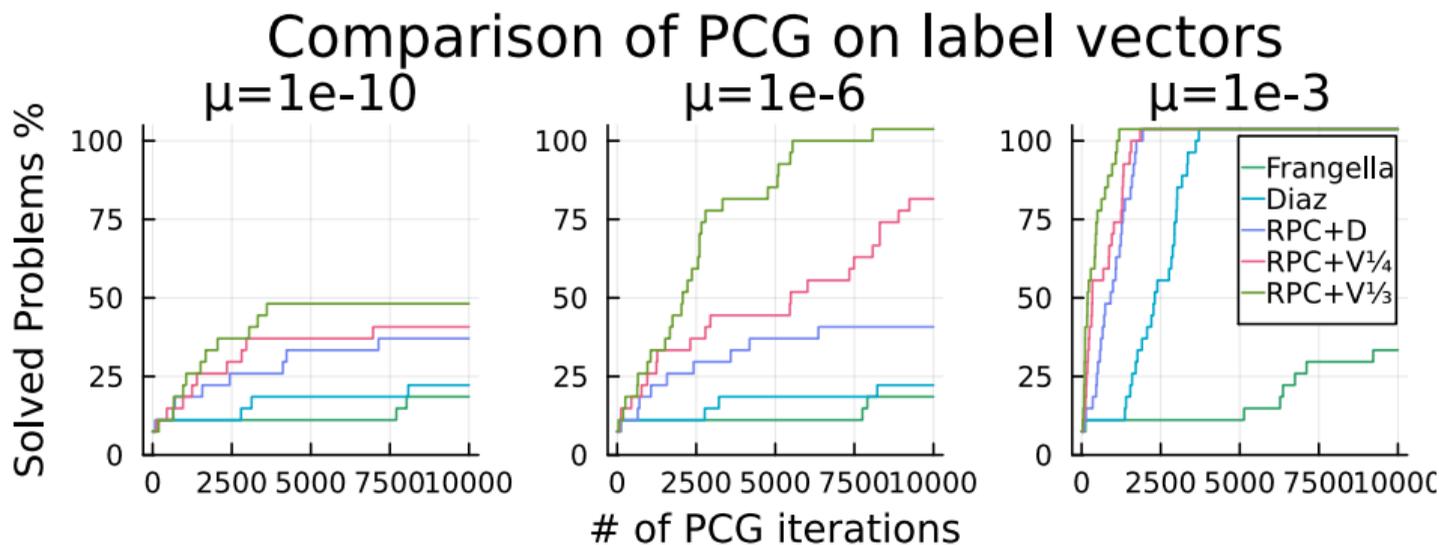
We standardized the predictors and formed the  $n \times n$  kernel matrix with entries

$$\mathbf{A}(i,j) = \exp\left(-\frac{\|\mathbf{x}_i - \mathbf{x}_j\|^2}{2d}\right) + \mu \delta(i,j), \quad \text{where } \mu \in \{10^{-3}, 10^{-6}, 10^{-10}\}.$$

We formed various matrix approximations  $\hat{\mathbf{A}}$  and tested performance as follows.

1. Use  $\hat{\mathbf{A}}$  as a preconditioner to solve  $\mathbf{Ax} = \mathbf{b}$ , where  $\mathbf{b}$  is the vector of labels.
2. Use  $\log \det(\hat{\mathbf{A}})$  as a direct estimator of  $\log(\det(\mathbf{A}))$ .
3. Use  $\hat{\mathbf{A}}$  as a preconditioner for stochastic log determinant estimation.

# Experimental results



**Figure:** We ran randomly pivoted Cholesky (RPC) with rank  $r = n^{1/2} = 141$  and used either (1) Frangella, Tropp, & Udell (2023); (2) Díaz, Epperly, Frangella, Tropp, & Webber (2023); (3) partial Cholesky + diagonal; (4) partial Cholesky + Vecchia with  $q = n^{1/4} = 11$  nonzeros in the residual sparsity pattern; or (5) partial Cholesky + Vecchia with  $q = n^{1/3} = 27$  nonzeros.

# Experimental results

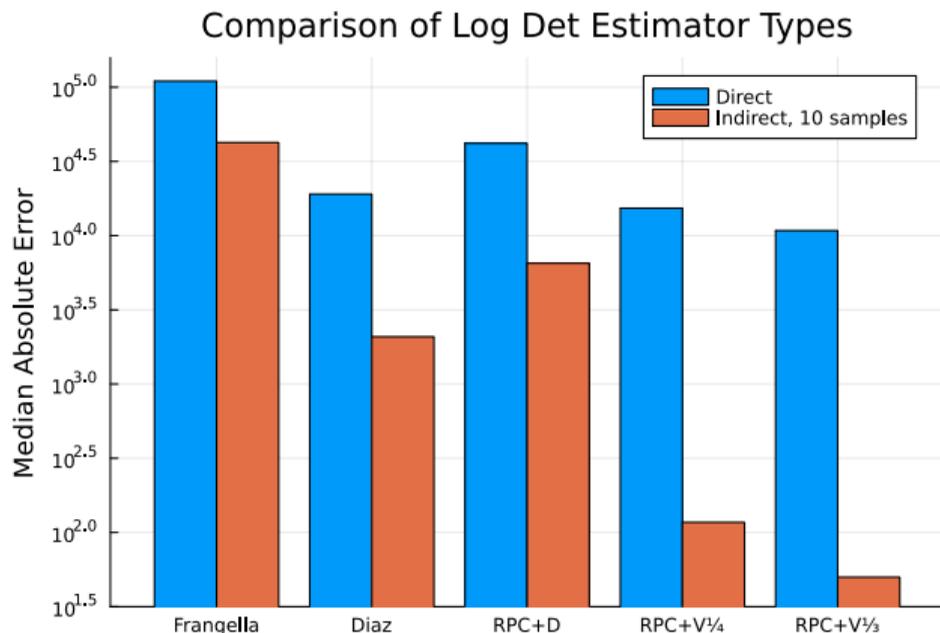


Figure: We set  $\mu = 10^{-3}$  and estimated  $\log(\det(\mathbf{A}))$  using (a) the direct estimator  $\log(\det(\hat{\mathbf{A}}))$  or (b) the stochastic log determinant estimator with  $\hat{\mathbf{A}}$  as a preconditioner.

# Takeaways

## Comparisons:

- For near-singular matrices, partial Cholesky + Vecchia provides the best approximation that's based on a partial Cholesky approximation.
- Partial Cholesky plus diagonal is okay but the Vecchia component helps a lot, especially with determinant estimation.

## Outlook:

- All approximations fail, but partial Cholesky + Vecchia fails more slowly as the eigenvalue lower bound  $\mu \downarrow 0$ .

# Vecchia optimality

Axelsson & Kaporin (1994, 2000) developed Vecchia optimality theory, and we generalized it to positive-semidefinite matrices.

## Definition (Kaporin condition number)

For any positive-semidefinite matrix  $\mathbf{A} \in \mathbb{C}^{n \times n}$  and any positive-semidefinite approximation  $\hat{\mathbf{A}} \in \mathbb{C}^{n \times n}$ , the Kaporin condition number is

$$\kappa_{\text{Kap}} = \frac{\left(\frac{1}{r} \text{tr}(\mathbf{A}\hat{\mathbf{A}}^+)\right)^r}{\text{vol}(\mathbf{A}\hat{\mathbf{A}}^+)}, \quad \text{where } r = \text{rank}(\mathbf{A}),$$

if  $\mathbf{A}$  and  $\hat{\mathbf{A}}$  share the same range. The Kaporin condition number is  $\kappa_{\text{Kap}} = \infty$  if  $\mathbf{A}$  and  $\hat{\mathbf{A}}$  have different ranges.

$\kappa_{\text{Kap}}$  is the average positive eigenvalue raised to the  $\text{rank}(\mathbf{A})$  power, divided by the product of the positive eigenvalues. By the AM-GM inequality,  $\kappa_{\text{Kap}} \geq 1$ .

# Vecchia optimality

## Theorem (Vecchia optimality)

For any positive-semidefinite matrix  $\mathbf{A} \in \mathbb{C}^{n \times n}$ , the Vecchia approximation  $\hat{\mathbf{A}} = \mathbf{P}\hat{\mathbf{C}}^{-1}\hat{\mathbf{D}}\hat{\mathbf{C}}^{-*}\mathbf{P}^*$  is the inverse Cholesky approximation with permutation  $\mathbf{P}$  and sparsity pattern  $\{S_i\}_{i=1}^n$  that achieves the smallest possible Kaporin condition number. When  $\hat{\mathbf{A}}$  and  $\mathbf{A}$  have the same range, the Kaporin condition number is

$$\kappa_{\text{Kap}} = \prod_{d_{\tilde{\mathbf{A}}}(\mathbf{e}_i, \text{span}\{\mathbf{e}_j\}_{j < i}) > 0} \frac{d_{\tilde{\mathbf{A}}}(\mathbf{e}_i, \text{span}\{\mathbf{e}_j\}_{j \in S_i})^2}{d_{\tilde{\mathbf{A}}}(\mathbf{e}_i, \text{span}\{\mathbf{e}_j\}_{j < i})^2},$$

where  $\tilde{\mathbf{A}} = \mathbf{P}^*\mathbf{A}\mathbf{P}$  is the permuted  $\mathbf{A}$  matrix.

# Implications of Vecchia optimality

$\kappa_{\text{Kap}}$  is important since it controls the error in linear algebra calculations.

Method	Error bound
Linear system, direct solver	$\frac{\ \hat{\mathbf{x}} - \mathbf{x}_*\ _{\mathbf{A}}^2}{\ \mathbf{x}_0 - \mathbf{x}_*\ _{\mathbf{A}}^2} \leq 2 \text{rank}(\mathbf{A}) \log(\kappa_{\text{Kap}})$
Linear system, iterative solver	$\frac{\ \mathbf{x}_t - \mathbf{x}_*\ _{\mathbf{A}}^2}{\ \mathbf{x}_0 - \mathbf{x}_*\ _{\mathbf{A}}^2} \leq \left[ \frac{3 \log(\kappa_{\text{Kap}})}{t} \right]^t$
Determinant, direct solver	$\log\left(\frac{\det \hat{\mathbf{A}}}{\det \mathbf{A}}\right) = \log(\kappa_{\text{Kap}})$
Determinant, iterative solver	$\mathbb{E} \left  \log\left(\frac{e^{st} \det \hat{\mathbf{A}}}{\det \mathbf{A}}\right) \right ^2 \leq \frac{4 \log(\kappa_{\text{Kap}})}{t}$

**Table:** Error bounds for direct and iterative solvers, assuming  $\text{tr}(\mathbf{A}\mathbf{A}^+) = \text{rank}(\mathbf{A})$ . Axelsson and Kaporin derived bound 2 for even  $t$  and we extended it. We derived bounds 1 and 4.

# Conclusions and prospects

## Structure of the Vecchia approximation

- + The Vecchia approximation is a superset of other approximations and has broad applicability.
- + We want to extend it to hierarchical matrices next.  
[Let me know if you're interested]

## Optimality theory

- + Vecchia is optimal for any given sparsity pattern and permutation, but finding the best sparsity pattern and permutation is NP-hard.
- + Our paper surveys simple heuristics for the sparsity pattern and permutation.
- + There's a lot more work to build up sophisticated sparsity patterns and prove recovery guarantees.

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Thank you for your attention! Does anyone have questions?