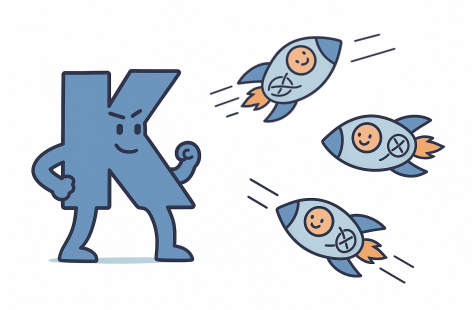


Diagrammatics for quantum symmetric pairs

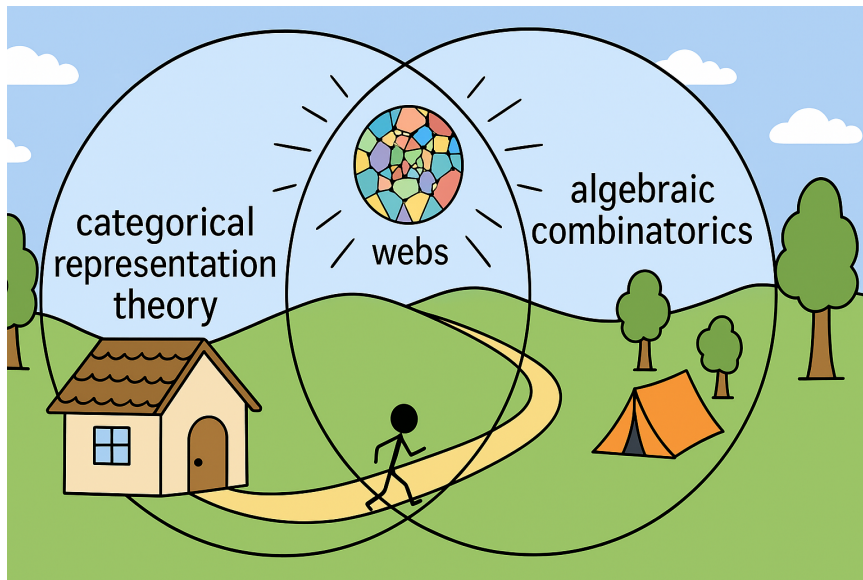


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Slides: alistairsavage.ca/talks, Paper: [arXiv:2507.12328](https://arxiv.org/abs/2507.12328)

The landscape



Outline

Goal: Use diagrammatic monoidal categories and module categories to study the representation theory of quantum symmetric pairs.

Overview:

- 1 Categorical background: Monoidal categories, string diagrams, and module categories
- 2 Quantum symmetric pairs
- 3 The disoriented skein category
- 4 The iquantum Brauer category
- 5 Connection to quantum symmetric pairs
- 6 Homework!

Strict linear monoidal categories

Fix a ground field $\mathbb{k}$.

A **strict linear monoidal category** is a $\mathbb{k}$ -linear category $\mathcal{C}$ equipped with

- a bifunctor (the **tensor product**) $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$, and
- a **unit object** $\mathbb{1}$,

such that, for objects A, B, C and morphisms f, g, h ,

- $(A \otimes B) \otimes C = A \otimes (B \otimes C)$,
- $\mathbb{1} \otimes A = A = A \otimes \mathbb{1}$,
- $(f \otimes g) \otimes h = f \otimes (g \otimes h)$,
- $1_{\mathbb{1}} \otimes f = f = f \otimes 1_{\mathbb{1}}$.

Remarks

- Every monoidal category is monoidally equivalent to a strict one.
- A **$\mathbb{k}$ -linear category** is one where the morphism spaces are $\mathbb{k}$ -modules, and composition and tensor product are $\mathbb{k}$ -bilinear.

String diagrams

Fix a strict monoidal category $\mathcal{C}$.

We will denote a morphism $f: A \rightarrow B$ by:



The **identity map** $1_A: A \rightarrow A$ is a string with no label:



We sometimes omit the object labels when they are clear or unimportant.

String diagrams

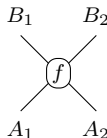
Composition is **vertical stacking** and tensor product is **horizontal juxtaposition**:

The diagram shows two equations. The first equation represents composition: a vertical line with a circle labeled f on top and a circle labeled g on the bottom is equal to a single vertical line with a circle labeled fg . The second equation represents the tensor product: a vertical line with a circle labeled f is tensored with a vertical line with a circle labeled g , which is equal to two vertical lines side-by-side, the left one with a circle labeled f and the right one with a circle labeled g .

The **interchange law** then becomes:

The diagram shows the interchange law: a vertical line with a circle labeled f is tensored with a vertical line with a circle labeled g , which is equal to a vertical line with a circle labeled f and a vertical line with a circle labeled g side-by-side, which is equal to a vertical line with a circle labeled f and a vertical line with a circle labeled g side-by-side, but with the g circle on top of the f circle.

A morphism $f: A_1 \otimes A_2 \rightarrow B_1 \otimes B_2$ can be depicted:



The oriented skein category

The **framed HOMFLYPT skein category**, or **oriented skein category** for short, $OS(q, t)$, is the $\mathbb{C}(q, t)$ -linear strict monoidal category generated by objects

$$\uparrow \quad \text{and} \quad \downarrow,$$

and morphisms

$$\begin{aligned} & \nearrow \nwarrow, \nwarrow \nearrow: \uparrow \otimes \uparrow \rightarrow \uparrow \otimes \uparrow, \\ \cup: \mathbb{1} \rightarrow \downarrow \otimes \uparrow, \quad \cap: \uparrow \otimes \downarrow \rightarrow \mathbb{1}, \quad \cup: \mathbb{1} \rightarrow \uparrow \otimes \downarrow, \quad \cap: \downarrow \otimes \uparrow \rightarrow \mathbb{1}, \end{aligned}$$

subject to the relations

$$\begin{aligned} \begin{array}{c} \nearrow \nwarrow \\ \nwarrow \nearrow \end{array} &= \begin{array}{c} \uparrow \uparrow \\ \downarrow \downarrow \end{array} = \begin{array}{c} \nwarrow \nearrow \\ \nearrow \nwarrow \end{array}, \quad \begin{array}{c} \nwarrow \nearrow \\ \nwarrow \nearrow \end{array} = \begin{array}{c} \nwarrow \nearrow \\ \nwarrow \nearrow \end{array}, \quad \begin{array}{c} \nwarrow \nearrow \\ \nwarrow \nearrow \end{array} = \begin{array}{c} \uparrow \downarrow \\ \downarrow \uparrow \end{array}, \quad \begin{array}{c} \nwarrow \nearrow \\ \nwarrow \nearrow \end{array} = \begin{array}{c} \downarrow \uparrow \\ \downarrow \uparrow \end{array}, \\ \nwarrow \nearrow - \nearrow \nwarrow &= (q - q^{-1}) \uparrow \uparrow, \quad \bigcirc \uparrow = t \uparrow = \uparrow \bigcirc, \quad \bigcirc = \frac{t - t^{-1}}{q - q^{-1}} \mathbb{1}, \\ \uparrow \cap &= \uparrow, \quad \downarrow \cup = \downarrow, \quad \text{where } \nwarrow \nearrow := \begin{array}{c} \uparrow \cap \\ \downarrow \cup \end{array}, \quad \nearrow \nwarrow := \begin{array}{c} \downarrow \cup \\ \uparrow \cap \end{array}. \end{aligned}$$

The oriented incarnation functor

Morphisms in $OS(q, t)$ are linear combinations of diagrams such as

$$: \downarrow \otimes \uparrow \otimes \downarrow \otimes \uparrow \otimes \uparrow \otimes \uparrow \otimes \downarrow \rightarrow \uparrow \otimes \downarrow \otimes \uparrow \otimes \uparrow \otimes \uparrow \otimes \downarrow \otimes \uparrow.$$

There is a **full** $\mathbb{C}(q)$ -linear monoidal functor

$$OS(q, q^m) \rightarrow U_q(\mathfrak{gl}(m))\text{-mod}, \quad \uparrow \mapsto V, \quad \downarrow \mapsto V^*,$$

where V is the quantum analogue of the natural module.

Connection to webs

The $U_q(\mathfrak{gl}_n)$ **web category** is a partial idempotent completion of $OS(q, q^n)$ (plus some additional relations).

Module categories

Suppose $\mathcal{C}$ is a strict monoidal category. A category $\mathcal{M}$ is a **right $\mathcal{C}$ -module category**, or **$\mathcal{C}$ -module** for short, if we have a functor

$$\mathcal{M} \times \mathcal{C} \rightarrow \mathcal{M}, \quad (M, C) \mapsto M \otimes C, \quad (f, g) \mapsto f \otimes g,$$

such that

$$\begin{aligned} M \otimes \mathbb{1} &= M, & f \otimes 1_{\mathbb{1}} &= f, \\ (M \otimes C) \otimes D &= M \otimes (C \otimes D), & (f \otimes g) \otimes h &= f \otimes (g \otimes h), \end{aligned}$$

for all objects M and morphisms f in $\mathcal{M}$, and objects C, D and morphisms g, h in $\mathcal{C}$.

There is a natural notion of a morphism of $\mathcal{C}$ -modules. These can be **strict** or **not strict**.

Diagrammatics for module categories

Suppose $\mathcal{C}$ is a monoidal category and $\mathcal{M}$ is a $\mathcal{C}$ -module.

We have a string diagram calculus for $\mathcal{M}$, except that we can only juxtapose on the **right** with morphisms in $\mathcal{C}$:

$$\begin{array}{c} | \\ \textcircled{f} \\ | \end{array} \otimes \begin{array}{c} | \\ \textcircled{g} \\ | \end{array} = \begin{array}{c} | \\ \textcircled{f} \textcircled{g} \\ | \end{array}, \quad f \in \text{Mor}(\mathcal{M}), \quad g \in \text{Mor}(\mathcal{C}).$$

We have a form of the interchange law:

$$\begin{array}{c} | \\ \textcircled{f} \\ | \end{array} \begin{array}{c} | \\ | \\ \textcircled{g} \\ | \end{array} = \begin{array}{c} | \\ \textcircled{f} \end{array} \begin{array}{c} | \\ | \\ \textcircled{g} \\ | \end{array} = \begin{array}{c} | \\ \textcircled{f} \end{array} \begin{array}{c} | \\ | \\ | \\ \textcircled{g} \\ | \end{array}, \quad f \in \text{Mor}(\mathcal{M}), \quad g \in \text{Mor}(\mathcal{C}).$$

Important:

- Morphisms in $\mathcal{M}$ can only appear on the **left** of a diagram.
- No horizontal juxtaposition of morphisms in $\mathcal{M}$.

Quantum symmetric pairs

Non-quantum case

Consider $\mathfrak{so}(m) \subseteq \mathfrak{gl}(m)$. Passing to universal enveloping algebras,

$U(\mathfrak{so}(m))$ is a Hopf subalgebra of $U(\mathfrak{gl}(m))$.

Thus, restriction

$$U(\mathfrak{gl}(m))\text{-mod} \rightarrow U(\mathfrak{so}(m))\text{-mod}$$

is a monoidal functor.

Quantum case

Things are not as nice in the quantum case.

Problem: $U_q(\mathfrak{so}(m))$ is not naturally a Hopf subalgebra of $U_q(\mathfrak{gl}(m))$.

The quantum enveloping algebra

Definition

The **quantum enveloping algebra** $U^{\imath} = U^{\imath}(\mathfrak{so}(m))$ is the $\mathbb{C}(q)$ -subalgebra of $U = U_q(\mathfrak{gl}(m))$ generated by

$$B_i = F_i - q^{-1}E_iK_i^{-1}, \quad 1 \leq i \leq m-1,$$

where E_i , F_i , and K_i are the usual generators of U .

Remarks

Taking $q \rightsquigarrow 1$ recovers $U(\mathfrak{so}(m)) \subseteq U(\mathfrak{gl}(m))$. So, $U^{\imath}$ is a quantum analogue of $U(\mathfrak{so}(m))$.

The quantum enveloping algebra

Important

$U^\imath$ is **not** a Hopf subalgebra of U . Instead, $U^\imath$ is a **coideal subalgebra**:

$$\Delta(U^\imath) \subseteq U^\imath \otimes U.$$

Thus, $U^\imath\text{-mod}$ is a **right U -module category**.

So, if $M, N \in U^\imath\text{-mod}$ and $X \in U\text{-mod}$, then

- we can form $M \otimes X \in U^\imath\text{-mod}$,
- we **cannot** form $M \otimes N$,
- we have the dual $X^* \in U\text{-mod}$,
- we do **not** have the dual M^* ,
- we have a trivial module $\text{triv} \in U\text{-mod}$,
- we have a trivial module $\text{Res}_{U^\imath}^U(\text{triv}) \in U^\imath\text{-mod}$.

Important motivating observation

Let $V \in U\text{-mod}$ be the quantum analogue of the natural module for $U = U_q(\mathfrak{gl}(m))$.

Then

$$V^* \not\cong V \quad \text{as } U\text{-modules,}$$

but

$$V^* \cong V \quad \text{as } U^{\iota}\text{-modules,}$$

as in the non-quantum setting.

So, to pass from $U\text{-mod}$ to $U^{\iota}\text{-mod}$, we need to add this isomorphism.

Oversimplified foreshadowing

To go from $U\text{-mod}$ to $U^{\iota}\text{-mod}$, we **only** need to add an isomorphism

$$\varphi: V^* \xrightarrow{\cong} V.$$

The disoriented skein category

The **disoriented skein category** $\mathcal{DS}(q, t)$ is the $OS(q, t)$ -module generated by the morphisms

$$\updownarrow : \downarrow \rightarrow \uparrow, \quad \star : \uparrow \rightarrow \downarrow,$$

which we call **toggles**, subject to the relations

$$\begin{array}{c} \uparrow \\ \circ \\ \star \\ \downarrow \end{array} = \uparrow, \quad \begin{array}{c} \downarrow \\ \star \\ \circ \\ \uparrow \end{array} = \downarrow, \quad \begin{array}{c} \uparrow \quad \uparrow \\ \diagdown \quad \diagup \\ \circ \quad \circ \\ \diagup \quad \diagdown \\ \downarrow \quad \downarrow \end{array} = \begin{array}{c} \uparrow \quad \uparrow \\ \diagdown \quad \diagup \\ \circ \quad \circ \\ \diagup \quad \diagdown \\ \downarrow \quad \downarrow \end{array}, \quad \begin{array}{c} \uparrow \\ \circ \\ \diagdown \quad \diagup \\ \downarrow \end{array} = q \begin{array}{c} \uparrow \\ \star \\ \downarrow \end{array}, \quad \begin{array}{c} \uparrow \\ \circ \\ \diagup \quad \diagdown \\ \downarrow \end{array} = q \begin{array}{c} \uparrow \quad \uparrow \\ \diagdown \quad \diagup \\ \downarrow \end{array}.$$

Guiding philosophy

The toggles impose the condition that $\uparrow$ and $\downarrow$ are isomorphic.

Warning!

$\mathcal{DS}(q, t)$ is **not** a monoidal category!

We don't have arbitrary juxtaposition. We can only juxtapose on the right with an $OS(q, t)$ diagram.

The disoriented incarnation functor

Recall that there is a full $\mathbb{C}(q)$ -linear monoidal functor

$$\mathbf{R}_{OS}: OS(q, q^m) \rightarrow U\text{-mod}, \quad \uparrow \mapsto V, \quad \downarrow \mapsto V^*,$$

and that we have an isomorphism

$$\varphi: V^* \xrightarrow{\cong} V \quad \text{of } U^l\text{-modules.}$$

The composition

$$U^l\text{-mod} \times OS(q, q^m) \xrightarrow{\text{id} \times \mathbf{R}_{OS}} U^l\text{-mod} \times U\text{-mod} \xrightarrow{\otimes} U^l\text{-mod}$$

makes $U^l\text{-mod}$ a $OS(q, q^m)$ -module.

Theorem (Salmasian–S.–Shen)

There is a **full** morphism of $OS(q, q^m)$ -modules

$$\begin{aligned} \mathbf{R}_{\mathcal{DS}}: \mathcal{DS}(q, q^m) &\rightarrow U^l\text{-mod}, \\ B_0 &\mapsto \text{Res}_{U^l}^U(\text{triv}), \quad \updownarrow \mapsto \varphi, \quad \nmid \mapsto \varphi^{-1}. \end{aligned}$$

The big picture (to be continued)

The following diagram commutes:

$$\begin{array}{ccc} \mathcal{DS} \times \mathcal{OS} & \xrightarrow{\quad \otimes \quad} & \mathcal{DS} \\ \searrow \mathbf{R}_{\mathcal{DS}} \times \mathbf{R}_{\mathcal{OS}} & & \swarrow \mathbf{R}_{\mathcal{DS}} \\ U^{\iota}\text{-mod} \times U\text{-mod} & \xrightarrow{\quad \otimes \quad} & U^{\iota}\text{-mod} \end{array}$$

Missing from the picture

Others have studied endomorphism algebras of U^{ι} -modules. How does that work fit into the above picture?

The quantum Brauer category

The **quantum Brauer category**, $\mathcal{B}(q, t)$, is a $\mathbb{C}(q, t)$ -linear category with objects B_r , $r \in \mathbb{N}$, and generating morphisms

$$\begin{aligned} \bigcup_r \mid : B_r &\rightarrow B_{r+2}, & \bigcap_r \mid : B_{r+2} &\rightarrow B_r, & r \in \mathbb{N}, \\ \mid_r \times \mid_s, \mid_r \times \mid_s &: B_{r+s+2} \rightarrow B_{r+s+2}, & r, s \in \mathbb{N}. \end{aligned}$$

The iquantum Brauer category

We impose the following relations on morphisms, for all $r, s \in \mathbb{N}$:

$$\left| \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array} \right|_r \quad \left| \begin{array}{c} \diagdown \diagup \\ \diagup \diagdown \end{array} \right|_s = \left| \right|_{r+s+2} = \left| \begin{array}{c} \diagdown \diagup \\ \diagdown \diagup \end{array} \right|_r \quad \left| \begin{array}{c} \diagup \diagdown \\ \diagup \diagdown \end{array} \right|_s = \left| \begin{array}{c} \diagup \diagdown \\ \diagup \diagdown \end{array} \right|_r \quad \left| \begin{array}{c} \diagdown \diagup \\ \diagdown \diagdown \end{array} \right|_s,$$

$$\left| \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array} \right|_r \quad \left| \begin{array}{c} \diagdown \diagup \\ \diagdown \diagup \end{array} \right|_s = (q - q^{-1}) \left| \right|_{r+s+2},$$

$$\left| \begin{array}{c} \cap \\ \cup \end{array} \right|_r = t \left| \right|_{r+1}, \quad \left| \begin{array}{c} \cap \\ \cap \end{array} \right|_r = q \left| \begin{array}{c} \cap \\ \cap \end{array} \right|_r, \quad \left| \begin{array}{c} \cup \\ \cup \end{array} \right|_r = q \left| \begin{array}{c} \cup \\ \cup \end{array} \right|_r, \quad \left| \begin{array}{c} \cap \\ \cup \end{array} \right|_r = \frac{t - t^{-1}}{q - q^{-1}} \left| \right|_r,$$

$$\left| \begin{array}{c} \cap \\ \cap \end{array} \right|_r = \left| \begin{array}{c} \cap \\ \cap \end{array} \right|_r, \quad \left| \begin{array}{c} \cap \\ \cap \end{array} \right|_r = \left| \begin{array}{c} \cap \\ \cap \end{array} \right|_r, \quad \left| \begin{array}{c} \cap \\ \cap \end{array} \right|_r = \left| \begin{array}{c} \cap \\ \cap \end{array} \right|_r,$$

(continued on next slide)

The quantum Brauer category

and

$$\begin{array}{c} | \\ \circlearrowleft f \\ | \end{array} \begin{array}{c} | \\ \circlearrowright g \\ | \end{array} := \begin{array}{c} | \\ \circlearrowleft f \\ | \end{array} \begin{array}{c} | \\ \circlearrowright g \\ | \end{array} = \begin{array}{c} | \\ \circlearrowleft f \\ | \end{array} \begin{array}{c} | \\ \circlearrowright g \\ | \end{array},$$

where f runs over **all** generating morphisms and g **only** runs over the generating morphisms

$$\begin{array}{c} | \\ \times \\ r \end{array} \begin{array}{c} | \\ \times \\ s \end{array} \text{ and } \begin{array}{c} | \\ \times \\ r \end{array} \begin{array}{c} | \\ \times \\ s \end{array}.$$

Remarks

- The quantum Brauer category is **not** a monoidal category.
- One cannot horizontally juxtapose diagrams or use the general interchange law for monoidal categories.
- Caps and cups can only appear on the left-hand side of diagrams.
- The category and its endomorphism algebras were developed by Molev, Wenzl, and Satori–Tubbenhauer. They are sometimes called the **q -Brauer category** and **q -Brauer algebras**.

The big picture (to be continued)

There is a $\mathbb{C}(q)$ -linear functor

$$\mathbf{R}_{\mathcal{B}}: \mathcal{B}(q, q^m) \rightarrow U^\iota\text{-mod}.$$

So now we have

$$\begin{array}{ccccc}
 \mathcal{DS} \times \mathcal{OS} & \xrightarrow{\quad \otimes \quad} & & \mathcal{DS} & \\
 \searrow \mathbf{R}_{\mathcal{DS}} \times \mathbf{R}_{\mathcal{OS}} & & & \swarrow \mathbf{R}_{\mathcal{DS}} & \\
 & U^\iota\text{-mod} \times U\text{-mod} & \xrightarrow{\quad \otimes \quad} & U^\iota\text{-mod} & \\
 \nearrow \mathbf{R}_{\mathcal{B}} \times \mathbf{R}_{\mathcal{OS}} & & & \nwarrow \mathbf{R}_{\mathcal{B}} & \\
 \mathcal{B} \times \mathcal{OS} & \xrightarrow{\quad \otimes \quad} & & \mathcal{B} &
 \end{array}$$

Missing from the picture

It is natural to expect that $\mathcal{B}(q, t)$ should be a right $\mathcal{OS}(q, t)$ -module category.

Module category structure on $\mathcal{B}(q, t)$ [Salmasian-S.-Shen]

We can make $\mathcal{B}(q, t)$ a $OS(q, t)$ -module as follows.

On **objects**, we define

$$B_r \otimes \uparrow := B_{r+1}, \quad B_r \otimes \downarrow := B_{r+1}.$$

On **morphisms**, we define

$$\begin{aligned} f \otimes \nearrow &:= f \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array}, & f \otimes \nwarrow &:= f \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array}, \\ f \otimes \uparrow &:= \begin{array}{c} f \\ \diagup \diagdown \\ \diagdown \diagup \end{array} = \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \\ f \end{array}, & f \otimes \downarrow &:= q^{-1}t \begin{array}{c} f \\ \diagup \diagdown \\ \diagdown \diagup \end{array} = q^{-1}t \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \\ f \end{array}, \\ f \otimes \downarrow &:= \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \\ f \end{array} = \begin{array}{c} f \\ \diagup \diagdown \\ \diagdown \diagup \end{array}, & f \otimes \uparrow &:= qt^{-1} \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \\ f \end{array} = qt^{-1} \begin{array}{c} f \\ \diagup \diagdown \\ \diagdown \diagup \end{array}. \end{aligned}$$

The big picture (to be continued)

So now we have

$$\begin{array}{ccc}
 \mathcal{DS} \times \mathcal{OS} & \xrightarrow{\quad \otimes \quad} & \mathcal{DS} \\
 \searrow \mathbf{R}_{\mathcal{DS} \times \mathcal{OS}} & & \swarrow \mathbf{R}_{\mathcal{DS}} \\
 & U^i\text{-mod} \times U\text{-mod} \xrightarrow{\quad \otimes \quad} U^i\text{-mod} & \\
 \nearrow \mathbf{R}_{\mathcal{B} \times \mathcal{OS}} & & \nwarrow \mathbf{R}_{\mathcal{B}} \\
 \mathcal{B} \times \mathcal{OS} & \xrightarrow{\quad \otimes \quad} & \mathcal{B}
 \end{array}$$

Remarks

- The top square commutes.
- The bottom square only commutes up to natural isomorphism.

Missing from the picture

What is the connection between $\mathcal{B}(q, t)$ and $\mathcal{DS}(q, t)$?

Functor from $\mathcal{DS}$ to $\mathcal{B}$

There is a **strict** morphism of $OS(q, t)$ -modules

$$\begin{aligned}\mathbf{F}: \mathcal{DS}(q, t) &\rightarrow \mathcal{B}(q, t), \\ \mathbf{F}(1) &= B_0, \quad \mathbf{F}\left(\begin{array}{c} \uparrow \\ \downarrow \end{array}\right) = 1_{B_1} = \mathbf{F}\left(\begin{array}{c} \downarrow \\ \uparrow \end{array}\right).\end{aligned}$$

Essentially, $\mathbf{F}$ forgets orientation and erases toggles.

One can explicitly define a (not strict!) morphism of $OS(q, t)$ -modules

$$\mathbf{G}: \mathcal{B}(q, t) \rightarrow \mathcal{DS}(q, t).$$

Theorem (Salmasian–S.–Shen)

The functors $\mathbf{F}$ and $\mathbf{G}$ are equivalences of $OS(q, t)$ -modules. More precisely,

$$\mathbf{FG} = \text{id}_{\mathcal{B}(q, t)}, \quad \mathbf{GF} \cong \text{id}_{\mathcal{DS}(q, t)}.$$

Thus, the **disoriented skein category** is equivalent to the **iquantum Brauer category** as right $OS(q, t)$ -module categories.

The big picture

Now we have

$$\begin{array}{ccccc}
 \mathcal{DS} \times \mathcal{OS} & \xrightarrow{\quad \otimes \quad} & & \mathcal{DS} & \\
 \uparrow \mathbf{G} \times \text{id} & \searrow \mathbf{R}_{\mathcal{DS}} \times \mathbf{R}_{\mathcal{OS}} & & \swarrow \mathbf{R}_{\mathcal{DS}} & \uparrow \mathbf{G} \\
 & U^{\mathcal{I}}\text{-mod} \times U\text{-mod} & \xrightarrow{\quad \otimes \quad} & U^{\mathcal{I}}\text{-mod} & \\
 \mathcal{B} \times \mathcal{OS} & \nearrow \mathbf{R}_{\mathcal{B}} \times \mathbf{R}_{\mathcal{OS}} & & \nwarrow \mathbf{R}_{\mathcal{B}} & \mathcal{B} \\
 & \xrightarrow{\quad \otimes \quad} & & &
 \end{array}$$

- $\mathbf{G}$ is an equivalence of categories.
- The left and right triangles commute ($\mathbf{R}_{\mathcal{B}} = \mathbf{R}_{\mathcal{DS}} \mathbf{G}$).
- The top rectangle commutes ($\mathbf{R}_{\mathcal{DS}}$ is strict).
- The bottom rectangle does **not** commute ($\mathbf{R}_{\mathcal{B}}$ is not strict).

Disoriented skein category versus iquantum Brauer category

The disoriented skein category $\mathcal{DS}$ has advantages over the iquantum Brauer category:

- commutativity of the diagram (incarnation functor is **strict** morphism of module categories),
- cups and caps in arbitrary positions,
- easier diagrammatics.

These advantages arise from the fact that $\mathcal{DS}$ contains the restrictions of the natural U -module **and** its dual, even though these are isomorphic as U^{ι} -modules.

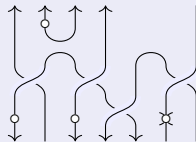
This gives extra flexibility, allowing the incarnation functor to be a **strict** morphism of $OS(q, t)$ -modules.

Other results

Basis theorems

One can describe explicit bases for the morphism spaces of $\mathcal{DS}(q, t)$, similar to the bases for Brauer, Kauffman, or oriented skein categories.

After defining toggles on arbitrary strands, the bases contain diagrams like:



As a corollary, one gets explicit bases for morphism spaces of $\mathcal{B}$ as well.

Symplectic and orthosymplectic

There are symplectic analogues of the above results. More generally, we can replace the orthogonal Lie algebra by the **orthosymplectic Lie algebra**.

Work in progress: Module subcategories

For each $d \in \mathbb{Z}$, we have incarnation functors

$$\mathbf{R}_{\mathcal{DS}}^{m,2n}: \mathcal{DS}(q, q^d) \rightarrow U^i(\mathfrak{osp}(m|2n))\text{-smod}, \quad m - 2n = d.$$

Conjecture

The kernels of the $\mathbf{R}_{\mathcal{DS}}^{m,2n}$, $m - 2n = d$, give a **complete** set of $OS(q, q^d)$ -submodules of $\mathcal{DS}(q, q^d)$.

These submodules form a chain:

$$\cdots \supseteq \ker \mathbf{R}_{\mathcal{DS}}^{m,2n} \supseteq \ker \mathbf{R}_{\mathcal{DS}}^{m+2,2n+2} \supseteq \cdots$$

Other work in progress

Diagrammatics for other quantum symmetric pairs.

Homework for spiders

Step 1: Take your favorite model, result, or theorem for webs (viewed as a monoidal category).

Step 2: Consider **module categories** over your webs, generated by an element corresponding to the **K -matrix** (the toggle).

Step 3: Profit. (Obtain web model, result, or theorem for **quantum symmetric pairs**.)

