

General Probabilistic Theories, tensor products, and projective transformations

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Abstract

Generalized Probabilistic Theories (GPTs) are theories of nature that have random features. A GPT must specify the set of states purporting to represent the physical reality, the allowable measurements, the rules for outcome statistics of the latter, and the composition rules for merging subsystems and creating a larger system. Examples include classical probability and quantum theory.

The composition rules alluded to above usually involve tensor products, including tensor products of normed spaces, convex sets and of cones. Among tensor products that have operational meaning in the GPT context, the projective and the injective product are the extreme ones, which leads to the natural question "How much do they differ?" considered already by Grothendieck and Pisier, resp. in the 1950s and 1980s.

We report on quantitative results concerning projective/injective discrepancy for finite-dimensional normed spaces. Some of the results are essentially optimal, but others can be likely improved. The methods involve a wide range of techniques from geometry of Banach spaces and random matrices. We also report on parallel results in the context of cones. Finally, we will encourage a more systematic study of convex bodies with the allowed morphisms being projective transformations.

- a few words about GPTs
- projective and injective tensor products and norms: definitions, notation
- historical background; the infinite dimensional case; qualitative vs. quantitative
- tensor products of cones vs. tensor products of normed spaces and convex bodies
- a selection of results and examples of tools from geometric functional analysis and random matrices

Buzzwords : CHSH inequality; Dvoretzky-Milman's theorem; p -summing norms; Chevet-Gordon's inequality; Grothendieck's inequalities; K -convexity & the MM^* -estimate

Generalized Probabilistic Theories or GPTs

A **general probabilistic theory** is a triple (V, C, u) , where: (i) V is a finite-dimensional real vector space; (ii) $C \subset V$ is a closed, convex, salient and generating cone; and (iii) u , called the **order unit** or the **unit effect**, is a functional in the interior of the dual cone

$$C^* := \{x^* \in V^* : x^*(x) \geq 0 \ \forall x \in C\}.$$

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Classical and quantum GPTs

In the **classical** case, $\Omega \subset \mathbb{R}^n$ is the set of (discrete) **probability densities**.

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Roughly speaking, **measurements** of a state give probabilistic outcomes according to the weights (p_j) and similarly in the quantum case.

This is **measurement “in a basis”**, other schemes are also allowed.

Definitions and notation : the projective norm

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Definitions and notation : duality and the injective norm

For the **smallest** “reasonable” norm on $X \otimes Y$ it is most convenient to appeal to **duality**: if $x^* \in X^*, y^* \in Y^*$, we want $x^* \otimes y^*$ to induce a functional on $X \otimes Y$ whose norm is $\|x^*\| \cdot \|y^*\|$, which implies

$$\|z\| \geq \max \{(x^* \otimes y^*)(z) : \|x^*\| \leq 1, \|y^*\| \leq 1\}. \quad (3)$$

Again, replacing “ \geq ” by “ $:=$ ” in (3) we get the definition of **injective** tensor norm $\|z\|_{X \otimes_{\varepsilon} Y}$ (or simply $\|z\|_{\varepsilon}$), denoted sometimes by $\|z\|_{X \check{\otimes} Y}$. Equivalently, $\|z\|_{\varepsilon}$ is the norm of z as a **bilinear form** on $X^* \times Y^*$.

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If the spaces are infinite dimensional, completions are required and there are reflexivity issues, but we will largely ignore this side of the story and – unless explicitly stated otherwise – will assume that $\dim X, \dim Y < \infty$.

An equivalent language: tensor products of convex sets

In geometric functional analysis, we often identify norms on a finite dimensional vector space V with symmetric **convex bodies**:

$$X = (V, \|\cdot\|) \rightarrow B_X := \{x : \|x\| \leq 1\} = \text{the unit ball of } X$$

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$$B_{X \otimes_{\pi} Y} := B_X \otimes_{\pi} B_Y \quad \text{and} \quad B_{X \otimes_{\varepsilon} Y} := (B_{X^*} \otimes_{\pi} B_{Y^*})^{\circ},$$

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One can likewise define projective and injective tensor products for not-necessarily-symmetric convex sets, most notably for **cones**.

Considering operators rather than tensors

Since $X^* \otimes Y$ is canonically isomorphic to $\mathcal{L}(X, Y)$, it is also possible to avoid talking about tensors and rephrase all questions in terms of **operators**. In that setting, if $z = \sum_i |y_i\rangle\langle x_i^*|$, then

$$\|z\|_\varepsilon = \|z : X \rightarrow Y\|,$$

the **operator norm**, while $\|z\|_\pi = \min \sum_i \|y_i\| \cdot \|x_i^*\|$ (the minimum over all representations) is the **nuclear norm**.

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$$\rho(X, Y) := \max_{z \in X \otimes Y, z \neq 0} \frac{\|z\|_\pi}{\|z\|_\varepsilon} = \max_{\|w: Y \rightarrow X^*\| \leq 1, \|z: X^* \rightarrow Y\| \leq 1} \operatorname{tr} wz.$$

The quantity $\rho(X, Y)$ quantifies **discrepancy** between $X \otimes_\pi Y$ and $X \otimes_\varepsilon Y$ and is arguably the most important concept of this presentation.

Tensor products of normed spaces were studied in detail by Grothendieck in 1950s. In particular, he proposed and studied 14 “natural tensor norms” and posed a number of open questions, one of which was whether the norms $\|\cdot\|_{X \otimes_{\pi} Y}$ and $\|\cdot\|_{X \otimes_{\varepsilon} Y}$ can be equivalent when $\dim X = \dim Y = \infty$.

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Also surprisingly, no quantitative analysis of the finite high-dimensional case was made until very recently. Such analysis is the main topic of this presentation. We also show that $\rho(X, Y) > 1$ in all nontrivial cases (i.e., $\min\{\dim W, \dim Y\} \geq 2$).

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Consider now any cone $C \subset \mathbb{R}^n$ and the classical cone $C' = \mathbb{R}_+^m \subset \mathbb{R}^m$.
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which means that each $z_k \in C$ (by the bipolar theorem for cones).

The equality case for cones

Consider now any cone $C \subset \mathbb{R}^n$ and the classical cone $C' = \mathbb{R}_+^m \subset \mathbb{R}^m$.

We want to show that

$$C \otimes_{\varepsilon} \mathbb{R}_+^m = C \otimes_{\pi} \mathbb{R}_+^m.$$

We have

$C \otimes_{\pi} \mathbb{R}_+^m = \text{conv}\{x \otimes y : x \in C, y \in \mathbb{R}_+^m\} = \text{conv}\{\sum_{k=1}^m x_k \otimes e_k : x_k \in C\}$
because each $y \in \mathbb{R}_+^m$ is a positive linear combination of e_k 's, and likewise

$$C^* \otimes_{\pi} \mathbb{R}_+^m = \text{conv}\{\sum_{k=1}^m y_k \otimes e_k : y_k \in C^*\}.$$

By definition, $\tau \in \mathbb{R}^n \otimes \mathbb{R}^m$ belongs to $C \otimes_{\varepsilon} \mathbb{R}_+^m$ iff

$$\langle \tau, \sigma \rangle \geq 0 \text{ for all } \sigma \in C^* \otimes_{\pi} \mathbb{R}_+^m.$$

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which means that each $z_k \in C$ (by the bipolar theorem for cones). That is, every $\tau \in C \otimes_{\varepsilon} \mathbb{R}_+^m$ belongs to $C \otimes_{\pi} \mathbb{R}_+^m$, so the two cones are equal.

The remarkable fact is that this is effectively the only case of equality.

Back to normed spaces: some special cases

If \mathcal{H}, \mathcal{K} are **Hilbert** (inner product) spaces, the situation is very simple:
 $\|\cdot\|_{\varepsilon}$ is the operator (spectral) norm, while $\|\cdot\|_{\pi}$ is the trace class norm
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is the Banach-Mazur distance. This allows to obtain some nontrivial information; for example using v, v^{-1} certifying $d(X, \mathcal{H}) \leq n^{1/2}$ as w, z in

$$\rho(X, \mathcal{H}) = \max_{\|w : \mathcal{H} \rightarrow X^*\| \leq 1, \|z : X^* \rightarrow \mathcal{H}\| \leq 1} \text{tr } wz$$

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we obtain $\rho(X, \mathcal{H}) \geq n^{1/2}$. The same circle of ideas allows to handle the case of different dimensions: $\rho(X, \mathcal{H}) \geq \min\{\dim X, \dim \mathcal{H}\}^{1/2}$.

Some special cases, cont'd

The same argument proves a **cute equality** $\rho(X, X^*) = \dim X$, but it doesn't help in the general case: by a 1981 result of Gluskin $\max\{d(E, F) : \dim E = \dim F = n\} = \Theta(n)$ and no nontrivial lower bound can be directly inferred.

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Here are other interesting special cases that can be handled. If (say) $\dim X \geq n$, then

$$\rho(X, \ell_1^n) \geq (n/2)^{1/2} \quad \text{and} \quad \rho(X, \ell_\infty^n) \geq (n/2)^{1/2}.$$

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The quantum case

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$$z = \sum_{j=1}^{d^2} \sum_{k=1}^{d^2} g_{jk} x_j \otimes y_k$$

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The quantum case, continued

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We note first that $z \in X \otimes_{\varepsilon} Y$ can be thought of as a **bilinear form** on $X^* \times Y^*$ and that $X^* = Y^*$ is (the self-adjoint part of) the **C^* -algebra** \mathcal{A} of $k \times k$ matrices with the usual operator norm. Thus we are in the realm of the Haagerup-Pisier **non-commutative Grothendieck inequality**, which says that for such bilinear form there are **states** φ, ψ on \mathcal{A} such that

$$|z(a, b)| \leq 2 \|z\|_{\varepsilon} \varphi(a^2)^{1/2} \psi(b^2)^{1/2} \quad \text{for all } a, b \in \text{Re } \mathcal{A}.$$

With this information, we need to upper-bound $\|z\|_{\pi}$, or the nuclear norm of $z : \mathcal{A} \rightarrow \mathcal{A}^*$.

The quantum case, conclusion

Again, we need to upper-bound the nuclear norm of $z : \mathcal{A} \rightarrow \mathcal{A}^*$, using

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$$z(a) = \sum_{i,j} \text{tr}(aE_{ij})z(E_{ij}), \quad \text{or } z = \sum_{i,j} |z(E_{ij})\rangle\langle E_{ij}|$$

where $E_{ij} = |u_i\rangle\langle u_j|$. For a single term, we have

$$\|z(E_{ij})\|_{\mathcal{A}^*} = \max_{\|b\|_{\mathcal{A}} \leq 1} |z(E_{ij}, b)| \leq 2\varphi(|E_{ij}|^2)^{1/2} \leq 2\lambda_i^{1/2}$$

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(note that $\psi(b^2) \leq 1$ if $\|b\|_{\mathcal{A}} \leq 1$) and summing over $1 \leq i, j \leq d$ gives $2d \sum_i \lambda_i^{1/2} \leq 2d^{3/2} = 2n^{3/4}$ as a bound on $\|z\|_\pi$ ($4n^{3/4}$ if we don't cheat).

The general case, or cases

Modulo logarithmic factors (indicated by $*$ in the Ω notation), we have:

- $X = Y, \dim X = n: \rho(X, X) = \Omega^*(n^{1/2})$ (almost optimal, see $X = \ell_1^n$)
- $\dim X = \dim Y = n: \rho(X, Y) = \Omega^*(n^{1/6})$
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The bounds can not be better than, respectively, $(2n)^{1/2}$, $n^{1/2}$, and again $n^{1/2}$. We know that $(2n)^{1/2}$ is not sharp, but we do not know whether the factor 2 can be removed. It is conceivable that all these quantities are actually $\Omega^*(n^{1/2})$ or even $\Omega(n^{1/2})$.

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While the above results are asymptotic, it is worthwhile noting that

- $\dim X, \dim Y \geq 2 \implies \rho(X, Y) \geq \frac{19}{18}$.

An Auerbach-type lemma

Just for fun, here is a simple lemma that is fundamental in proving this 2-dimensional lower bound and which is susceptible to an intuitive geometrical interpretation:

Let X be a Banach space of dimension at least 2. Then there exist vectors $e_1, e_2 \in X$, $e_1^, e_2^* \in X^*$ such that for any $i, j \in \{1, 2\}$ we have $\|e_i\|_X = \|e_j^*\|_{X^*} = 1$, $e_j^*(e_i) = \delta_{ij}$, and moreover*

$$\|e_1 + e_2\|_X \leq 3/2.$$

Once established, this allows to mimic the entanglement-detecting CHSH-inequality to obtain a non-trivial $19/18$ lower bound for $\rho(X, Y)$. This is the bound shown in our paper; we have a more complicated argument that gives $\frac{8}{7}$, and we suspect that the right number is $\sqrt{2}$.

Affine vs. projective transformations

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One illustration of this difference is the fact that all convex quadrangles are projectively equivalent, while obviously they are not all affinely equivalent. This suggests investigating the class of n -dimensional (not necessarily-symmetric) convex bodies by identifying bodies that are projectively equivalent, and with some projective version of the Banach-Mazur distance quantifying their similarity.

The toolbox for the general case

This is again based on $\rho(X, Y) = \max_{\|w: Y \rightarrow X^*\| \leq 1, \|z: X^* \rightarrow Y\| \leq 1} \operatorname{tr} wz$ and an appropriate relaxation of the choices $w = v, z = v^{-1}$. First, we define the **factorization constant** of X through Y as

$$f(X, Y) := \inf_{u, v} \{ \|u : X \rightarrow Y\| \cdot \|v : Y \rightarrow X\| : vu = \operatorname{Id}_X \},$$

which allows $\dim X \neq \dim Y$ and means that a subspace “**well-isomorphic**” to X is “**well-complemented**” in Y .

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$$\text{wf}(X, Y) := \inf_{u, v} \{ \mathbf{E} [\|u : X \rightarrow Y\| \cdot \|v : Y \rightarrow X\|] : \mathbf{E} [vu] = \text{Id}_X \},$$

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Clearly $\text{wf}(X, Y) \leq f(X, Y) \leq d(X, Y)$ and one easily checks that

$$\rho(X', Y) \leq \text{wf}(X', X)\rho(X, Y) \quad \text{and} \quad \rho(X, Y) \geq \frac{\dim X}{\text{wf}(X, Y^*)}.$$

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$$\|\text{Id} : \ell_2^n \rightarrow X\| \cdot \mathbf{E}\|g\|_Y + \|\text{Id} : \ell_2^n \rightarrow Y\| \cdot \mathbf{E}\|g\|_X,$$

where g is the **standard Gaussian vector** on \mathbb{R}^n .

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For an appropriate representation of X of \mathbb{R}^n , the first two factors give $d(X, \ell_2^n) \leq n^{1/2}$. The last two factors are essentially the same as **spherical means**, or **mean widths**, which can be controlled by the **MM*-estimate**.

Mean (half-)width of $K \subset \mathbb{R}^n$ and the MM^* -estimate

If $|u| = 1$ and $w(K, u) := \sup_{x \in K} \langle u, x \rangle = \|u\|_{K^\circ}$, then $w(K, u) + w(K, -u)$ is the width of K in the direction of u . The average over u is the mean **half-width** of K .

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The MM^* -estimate says that, for some **well-balanced linear image** \tilde{K} of any centrally symmetric convex body $K \subset \mathbb{R}^n$ we can achieve $w(\tilde{K}) \cdot w(\tilde{K}^\circ) = O(\log n)$.

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Some additional tweaking is needed since we need to reconcile two requirements for the representation of X of \mathbb{R}^n , the one witnessing $d(X, \ell_2^n)$ and the other consistent with the MM^* -estimate, but ultimately gathering all bounds we get

$$\rho(X, X) = \Omega\left(\frac{\dim X}{d(X, \ell_2^n) \log^3 n}\right) \geq \Omega\left(\frac{n^{1/2}}{\log^3 n}\right),$$

as needed.

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Let X be a normed space of dimension n . Then for every $1 \leq A \leq n^{1/2}$ at least one of the following holds

- 1 X contains a subspace E of dimension $d = \Omega(n^{1/2})$ such that $d(E, \ell_\infty^d) = O(A\sqrt{\log n})$.
- 2 X^* contains a subspace F of dimension $d = \Omega(n^{1/2})$ such that $d(F, \ell_\infty^d) = O(A\sqrt{\log n})$.
- 3 X contains a subspace H of dimension $d = \Omega(A^2/\log n)$ such that $d(H, \ell_2^d) \leq 4$ and, additionally, H is $O(\log n)$ -complemented in X .

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Since subspaces λ -isomorphic to ℓ_∞^d are automatically λ -complemented, each of the conditions above leads to an upper bound on $\text{wf}(\ell_p^d, X)$ for the appropriate $p \in \{1, 2, \infty\}$. Given that $\rho(\ell_p^d, \ell_{p'}^{d'})$ are known, every combination of these conditions for X and Y leads to a lower bound on $\rho(X, Y)$, and the final step is optimizing over A .

THANK YOU!