The G-Shi Arrangement: Games on Paths, Trees, and More

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Background	Three Rows Game	Repetitions	Generalizing to Trees	Beyond Trees	Conclusion
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- \circ The G-Shi arrangement
- The Pak-Stanley algorithm

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- $\circ~$ The $G\text{-}\mathsf{Shi}$ arrangement
- The Pak-Stanley algorithm
- 2 The Three Rows Game
 - \circ Path graphs
 - $\circ\,$ The Path Repetition Theorem

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 - $\circ~$ The Path Repetition Theorem
- **3** The Three Rows Game for Trees
 - \circ Acyclicity
 - $\circ~$ Star graphs

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- **3** The Three Rows Game for Trees
 - Acyclicity
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- 4 The Three Rows Game for All Graphs
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BACKGROUND

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Hyperplanes



A hyperplane in \mathbb{R}^n is an affine subspace of dimension n-1.

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Hyperplane Arrangements



A *hyperplane arrangement* is a finite collection of hyperplanes. The hyperplane arrangement separates the vector space into disjoint *chambers* or *regions*.

G-Shi Arrangement

Given a graph G, the G-Shi arrangement includes the hyperplanes $x_i - x_j = 0$ and $x_i - x_j = 1$ for each edge $\{i, j\}$ where i < j.



If G has n vertices and m edges, the G-Shi arrangement exists in n dimensions and has 2m hyperplanes.

G-Parking Functions

Let G be a graph on n + 1 vertices: $\{0, 1, \dots, n - 1\}$ and a sink q.

Definition

A *G*-parking function is an *n*-tuple $(a_0, a_1, ..., a_{n-1})$ such that for any non-empty subset $S \subseteq \{0, ..., n-1\}$, there exists $v \in S$ such that $a_v < \operatorname{outdeg}_S(v)$.

This is also known as a *superstable configuration* in chip-firing.

Three Rows Game Background Generalizing to Trees

An Example of *G*-Parking Functions



{0}	{1}	{2}	$\{0,1\}$	$\{0, 2\}$	$\{1, 2\}$	$\{0, 1, 2\}$
			$a_0 < 1$	$a_0 < 2$	$a_1 < 2$	$a_0 < 1$ or
$a_0 < 2$	$a_1 < 3$	$a_2 < 2$	or	or	or	$a_1 < 1$ or
			$a_1 < 2$	$a_2 < 2$	$a_2 < 1$	$a_2 < 1$

These constraints yield the G-parking functions (0, 0, 0), (0, 0, 1), (0, 1, 0), (1, 0, 0), (0, 1, 1), (1, 0, 1), (1, 1, 0), and (0, 2, 0)

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An Example of G-Parking Functions



{0}	$\{1\}$	$\{2\}$	$\{0, 1\}$	$\{0, 2\}$	$\{1, 2\}$	$\{0, 1, 2\}$
			$a_0 < 1$	$a_0 < 2$	$a_1 < 2$	$a_0 < 1$ or
$a_0 < 2$	$a_1 < 3$	$a_2 < 2$	or	or	or	$a_1 < 1$ or
			$a_1 < 2$	$a_2 < 2$	$a_2 < 1$	$a_2 < 1$

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An Example of G-Parking Functions



{0}	{1}	$\{2\}$	$\{0, 1\}$	$\{0, 2\}$	$\{1, 2\}$	$\{0, 1, 2\}$
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An Example of G-Parking Functions



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			$a_0 < 1$	$a_0 < 2$	$a_1 < 2$	$a_0 < 1$ or
$a_0 < 2$	$a_1 < 3$	$a_2 < 2$	or	or	or	$a_1 < 1$ or
			$a_1 < 2$	$a_2 < 2$	$a_2 < 1$	$a_2 < 1$

These constraints yield the G-parking functions (0,0,0), (0,0,1), (0,1,0), (1,0,0), (0,1,1), (1,0,1), (1,1,0), and (0,2,0)

Appending Sinks: G_{\bullet}

Let G = (V, E) be a graph. Then, G_{\bullet} is the graph obtained from G by adding a vertex q, which we call the *sink*, and an edge between q and each $v \in V$. For example, let $G = P_3$.



Then, G_{\bullet} is the graph below.



Shi Adjacency Graph

The Shi adjacency graph is defined from the G-Shi arrangement by letting each region be a vertex and placing an edge between bordering regions.



Directing the Shi Adjacency Graph

First, we label the region given by $x_0 > x_1 > \cdots > x_{n-1}$ and $x_0 - x_{n-1} < 1$ by R_0 , and call it the *base region*. Then, we direct the edges away from R_0 .



Above is the example for $G = K_3$.

Pak-Stanley Labels

In the Pak-Stanley Algorithm, we label each edge $x_i - x_j = 0$ with i and each edge $x_i - x_j = 1$ with j. Then, starting from R_0 , we label the vertices by increasing coordinate i along edges labeled i.



Pak-Stanley Labels

The set of Pak-Stanley labels for the regions of the G-Shi arrangement are called *the Pak-Stanley labels for* G; these are the same as the G_{\bullet} -parking functions (Hopkins-Perkinson).

Theorem (Hopkins-Perkinson, Corollary 2.8)

Every G_{\bullet} -parking function occurs as a label in the Pak-Stanley algorithm on the G-Shi arrangement.

Furthermore, for complete graphs, each $(K_n)_{\bullet}$ -parking function appears exactly once. However, this is not the case for graphs which not are not complete.



An Example: Path Graphs

We started by studying a simple family of graphs: path graphs.



The Shi adjacency digraph of P_1 , P_2 , P_3 , and P_4 :



These graphs look like subdivided (n-1)-dimensional cubes.

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Repetitions in the Labels for Path Graphs

The Pak-Stanley labels on the Shi adjacency graph for P_3 .



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THREE ROWS GAME



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$x_0 - x_1 < 0$	
$0 < x_0 - x_1 < 1$	
$x_0 - x_1 > 1$	

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Three Rows Game on Paths

The Three Rows Game for P_n uses the following board:

0	1	2	 n-3	n-2
1	2	3	 n-2	n-1

To play, pick one cell in each column and keep track of the values selected.

Example

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Histories

The vertices in the Shi adjacency digraph correspond to *histories* in the Three Rows Game.

Example

What history corresponds to this vertex?



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Histories

The vertices in the Shi adjacency digraph correspond to *histories* in the Three Rows Game.

Example

What history corresponds to this vertex?



Outcomes

The *outcomes* of the Three Rows Game correspond to the Pak-Stanley labels.

Example

What is the Pak-Stanley label for this vertex?



Outcomes

Question

Given an outcome of the Three Rows Game (or Pak-Stanley label), can we determine which vertex it came from?

Example

Which vertex does the outcome $\{1\}$ correspond to?



Outcomes

Question

Given an outcome of the Three Rows Game (or Pak-Stanley label), can we determine which vertex it came from?

Example

Which vertex does the outcome $\{1\}$ correspond to?



There are multiple solutions! This corresponds to the fact that the label $\left(0,1,0\right)$ appears twice, as discussed previously.

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REPETITIONS

Counting Repetitions

Question

For a specific G-parking function, how many times will it appear in the P_n -Shi arrangement?

0	1	2	3
1	2	3	4

How many histories yield the following outcomes:

 $\{0, 2, 2\}$?
Counting Repetitions

Question

For a specific $G\mbox{-}\mathsf{parking}$ function, how many times will it appear in the $P_n\mbox{-}\mathsf{Shi}$ arrangement?

0	1	2	3	
1	2	3	4	

How many histories yield the following outcomes:

$$\{0,2,2\}? \longrightarrow 1$$

0	1	2	3	
1	2	3	4	

Counting Repetitions

Question

For a specific G-parking function, how many times will it appear in the P_n -Shi arrangement?

0	1	2	3
1	2	3	4

How many histories yield the following outcomes:

 $\{1, 2, 4\}$?

Counting Repetitions

Question

For a specific G-parking function, how many times will it appear in the P_n -Shi arrangement?

0	1	2	3	
1	2	3	4	

How many histories yield the following outcomes:

$$\{1,2,4\}? \longrightarrow 3$$

0	1	2	3	0	1	2	3	0	1	2	3
1	2	3	4	1	2	3	4	1	2	3	4

Path Repetition Theorem

Definition

A run of length n is a subsequence of the Pak-Stanley label of the form $0,1,1,\ldots,1,1,0$ where there are n 1's.

Theorem (C. Bennett, A. Mock, R. Truax)

Let **p** be a Pak-Stanley label on P_n . If the length of a run r is denoted l(r), then the number of vertices in $\Gamma \mathscr{S}(P_n)$ with label **p** is

$$\prod_{l \neq r \text{ in } \mathbf{p}} (l(r) + 1).$$

runs r in p

Question

How many times will the Pak-Stanley label (0, 1, 1, 1, 0, 2, 0) appear in the P_7 -Shi arrangement?

Question

How many times will the Pak-Stanley label $\left(0,1,1,1,0,2,0\right)$ appear in the $P_{7}\text{-Shi}$ arrangement?

Solution

There is one run of length 3:

 $(0, \mathbf{1}, \mathbf{1}, \mathbf{1}, 0, 2, 0)$

so the label appears 3 + 1 = 4 times.

Question

How many times will the Pak-Stanley label (0, 1, 1, 0, 1, 1, 0) appear in the P_7 -Shi arrangement?

Question

How many times will the Pak-Stanley label (0, 1, 1, 0, 1, 1, 0) appear in the P_7 -Shi arrangement?

Solution

There are two runs of length 2:

 $(0, \mathbf{1}, \mathbf{1}, 0, \mathbf{1}, \mathbf{1}, 0)$

so the label appears (2+1)(2+1) = 9 times.

Results from the Three Rows Game

Fact (C. Bennett, A. Mock, R. Truax)

The number of regions in the P_n -Shi arrangement is equal to 3^{n-1} .

Explanation

Recall that each region corresponds to a history in the Three Rows Game. For the P_n -Shi arrangement, there are n-1 columns.

0	1	•••	n-2
1	2		n-1

We choose one cell in each column, so there are 3^{n-1} histories.

Results from the Three Rows Game

Fact (C. Bennett, A. Mock, R. Truax)

The number of sinks in the Shi adjacency digraph of P_n is equal to 2^{n-1} .

Explanation

A *sink* is a vertex in the Shi adjacency digraph with outdegree 0. For paths (and indeed trees), this corresponds to having no blanks in the history:

There are n-1 columns, and two choices for each. Thus, there are 2^{n-1} such histories.

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The *T*-Three Rows Game

The T-Three Rows Game has a colum

$$\begin{array}{c|c}
i \\
j \\
\hline j
\end{array}$$
 for each edge $\{i, j\}$.



0	1	1	3	
1	2	3	4	

Acyclicity and the Analogy

Acyclicity allows us to choose any sequence of moves.

 $\,\hookrightarrow\,$ a cycle could introduce a contradiction.



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Since trees have no cycles, this issue doesn't arise. Thus, the Shi adjacency digraph still looks like a barycentrically subdivided (n-1)-cube, it has 3^{n-1} vertices and 2^{n-1} sinks, etc.

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An Example: The Star Graph



0	0	0	···· 0		0
1	2	3	• • •	n-1	n

Generalizing to Trees

An Example: The Star Graph



0	0	0	··· 0		0
1	2	3	• • •	n-1	n

How many histories of the game give the outcome $\{1, 2, 4\}$?

0	0	0	0	0	0	0
1	2	3	4	5	6	7

Generalizing to Trees

An Example: The Star Graph



0	0	0	• • •	0	0
1	2	3	• • •	n-1	n

How many histories of the game give the outcome $\{1, 2, 4\}$?

0	0	0	0	0	0	0
1	2	3	4	5	6	7

Generalizing to Trees

An Example: The Star Graph



0	0	0	• • •	0	0
1	2	3	• • •	n-1	n

How many histories of the game give the outcome $\{0, 0, 1, 4, 5\}$?

0	0	0	0	0	0	0
1	2	3	4	5	6	7

Generalizing to Trees

An Example: The Star Graph



0	0	0	• • •	0	
1	2	3	• • •	n-1	n

How many histories of the game give the outcome $\{0, 0, 1, 4, 5\}$?

0	0	0	0	0	0	0
1	2	3	4	5	6	7

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An Example: The Star Graph

0	0	0		0	0
1	2	3	• • •	n-1	n

How many histories of the game give the outcome $\{0, 0, 1, 4, 5\}$?

0	0	0	0	0	0	0
1	2	3	4	5	6	7

$$\binom{4}{2} = 6$$

Classifying Repetitions in the Star Graph

Theorem (C. Bennett, A. Mock, R. Truax)

The number of times an outcome repeats in the S_n -Three Rows Game is

$$\begin{pmatrix} c_0 + c_{\Box} \\ c_0 \end{pmatrix} \cdot$$
of 0s # of \Box s

Corollary

The number of times a Pak-Stanley label $\mathbf{p} = (p_0, \dots, p_n)$ repeats in the S_n -Shi arrangement is

$$\binom{n-p_1-\cdots-p_n}{p_0}.$$

Notice that any outcome which has no \Box s appears precisely once.

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Cycle Graphs

If G has a cycle, some histories of the G-Three Rows Game are *illegal* (they correspond to nonexistent regions). The simplest case is when G is a cycle: the cycle graph C_n :



Here are two examples of illegal histories for $G = C_n$:

0	1	• • •	n-2	0		0	1	•••	n-2	0
					and					
1	2		n-1	n-1		1	2		n-1	n-1

Here are two examples of illegal histories for $G = C_n$:

0	1	•••	n-2	0		0	1	•••	n-2	0
					and					
1	2		n-1	n-1		1	2	•••	n-1	n-1

To see why these histories are illegal, notice that the first n-1 choices ("path choices") become the inequalities

$$x_0 - x_1 < 0$$
 $x_1 - x_2 < 0$ \cdots $x_{n-2} - x_{n-1} < 0$

Here are two examples of illegal histories for $G = C_n$:

0	1	•••	n-2	0		0	1	•••	n-2	0
					and					
1	2		n-1	n-1		1	2	• • •	n-1	n-1

To see why these histories are illegal, notice that the first n-1 choices ("path choices") become the inequalities

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Here are two examples of illegal histories for $G = C_n$:

0	1	• • •	n-2	0		0	1	•••	n-2	0
					and					
1	2		n-1	n-1		1	2	• • •	n-1	n-1

To see why these histories are illegal, notice that the first n-1 choices ("path choices") become the inequalities

$$x_0 - x_{n-1} < 0$$

But the final choice ("cycle choice") becomes $0 < x_0 - x_{n-1} < 1$ and $x_0 - x_{n-1} > 1$ respectively (a contradiction).

A Classification of Illegal Histories

Case 1: All the path choices are T or M.

This case has n + 1 illegal histories:



Case 2: All the path choices are M or B.

This case has $2^n - 1$ illegal histories: all the path choices are M or B and the cycle choice is M or T, except the history of all Ms.

Corollaries of the Classification

The classification yields the following results:

Theorem (C. Bennett, A. Mock, R. Truax) There are $2^n + n$ illegal histories in the C_n -Three Rows Game.

Corollary

There are $3^n - 2^n - n$ legal histories in the C_n -Three Rows Game.

Corollary

There are $3^n - 2^n - n$ regions in the C_n -Shi arrangement.

It also gives a Cycle Repetition Theorem analogous to the Path Repetition Theorem (but more complicated).

Maximal Labels

A maximal G_{\bullet} -parking function is one that cannot be made any larger and remain a G_{\bullet} -parking function. For example, (1,1,0,1,1) is a maximal $(P_5)_{\bullet}$ -parking function.

Proposition

Suppose G has m edges. Then any maximal G_{\bullet} -parking function has coordinate sum m; conversely, any G_{\bullet} -parking function with coordinate sum m is maximal.

Corollary

An outcome of the G-Three Rows Game corresponds to a maximal G_{\bullet} -parking function if and only if the middle row is never used.



Uniqueness of Maximal Labels

Theorem (C. Bennett, A. Mock, R. Truax)

Any outcome of the *G*-Three Rows Game which never uses the middle row is the result of a unique legal history.

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Any outcome of the *G*-Three Rows Game which never uses the middle row is the result of a unique legal history.

Proof.

The idea is to choose a legal history ${\bf h}$ inducing a maximal outcome ${\bf o},$ and show that any other history ${\bf h}'$ inducing ${\bf o}$ is illegal.

Indeed, if \mathbf{h}' is not illegal, then it would need to have infinitely many columns distinct from \mathbf{h} , which is impossible.

Uniqueness of Maximal Labels

Theorem (C. Bennett, A. Mock, R. Truax)

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The idea is to choose a legal history ${\bf h}$ inducing a maximal outcome ${\bf o},$ and show that any other history ${\bf h}'$ inducing ${\bf o}$ is illegal.

Indeed, if \mathbf{h}' is not illegal, then it would need to have infinitely many columns distinct from \mathbf{h} , which is impossible.

Corollary

Any maximal G_•-parking function appears uniquely.

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CONCLUSION

Open Problems: Counting Regions

Question

How many regions are in the G-Shi arrangement of any graph?

Example

The G-Shi arrangements of these graphs have different numbers of regions.


Open Problems: Counting Repetitions

Question

How many times does a given label repeat in the $T\mbox{-}{\rm Shi}$ arrangement of a tree graph T?

- We can find the answer using breadth-first search.
- Does a faster algorithm exist?

Question

How many times does a given label repeat in the G-Shi arrangement for a general graph G?

Open Problems: Sinks

Every vertex with a maximal Pak-Stanley label is a sink in the Shi adjacency digraph.

Question

Does every sink in the Shi adjacency digraph have a maximal Pak-Stanley label?

• We have shown that this is true for tree graphs.

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THANK YOU!

Appendix •0000

Appendix

Uniqueness in the *T*-Three Rows Game

Theorem (C. Bennett, A. Mock, R. Truax)

Any outcome of the *T*-Three Rows Game with no $\Box s$ (i.e. with *n* numbers) is the result of a unique history.

Uniqueness in the *T*-Three Rows Game

Theorem (C. Bennett, A. Mock, R. Truax)

Any outcome of the *T*-Three Rows Game with no $\Box s$ (i.e. with *n* numbers) is the result of a unique history.

Proof.

First, we introduce the concept of *distance* between outcomes:

$$d(\mathbf{o}, \mathbf{o}') = \sum_{i=0}^{n} |\mathbf{m}_{\mathbf{o}}(i) - \mathbf{m}_{\mathbf{o}'}(i)|.$$
occurences of *i* in o
occurences of *i* in o'

Uniqueness in the *T*-Three Rows Game

Theorem (C. Bennett, A. Mock, R. Truax)

Any outcome of the T-Three Rows Game with no $\Box s$ (i.e. with n numbers) is the result of a unique history.

Proof.

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$$d(\mathbf{o}, \mathbf{o}') = \sum_{i=0}^{n} |\mathbf{m}_{\mathbf{o}}(i) - \mathbf{m}_{\mathbf{o}'}(i)|.$$
occurences of *i* in o occurences of *i* in o'

Since T is a tree, we can relabel the vertices so that every column (after the first) contains one old vertex and one new vertex.

Uniqueness in the *T*-Three Rows Game (Cont.)

Proof (Cont).

Suppose that h_1 and h_2 give the same outcome o. Since our outcome has no \Box s, h_1 can be transformed into h_2 using "swaps":



Uniqueness in the *T*-Three Rows Game (Cont.)

Proof (Cont).

Suppose that h_1 and h_2 give the same outcome o. Since our outcome has no \Box s, h_1 can be transformed into h_2 using "swaps":



Now, the first swap increases the distance by 2. However, each swap afterwards doesn't decrease the distance (since there is one new and one old vertex). This gives a contradiction.

Cycle Graphs

Remember that if there is a cycle, some histories are *illegal*. In the case $G = C_n$, we can count these illegal histories:



3 Any history where the first n-1 choices are all M or B and the last choice is M or T (except the history of all Ms).

Thus, the number of regions in the C_n -Shi arrangement is

$$3^{n} - 2 - (n - 1) - (2^{n} - 1) = 3^{n} - 2^{n} - n.$$

Patterns

A *pattern* is a maximal sequence of weakly-increasing choices with exactly one middle choice.

These histories contain patterns:

0	1	2	3	4
1	2	3	4	5

0	1	2	3	4
1	2	3	4	5

These histories contain no nontrivial patterns:

0	1	2	3	4
1	2	3	4	5

0	1	2	3	4
1	2	3	4	5